### Int.tract.iw, Repair-lkd Planning and Scheduling for Shuttle Payload Operations

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Abstract This paper describes the 1 DATA-CHASER Automated Planner/Scheduler (1 DCAPS) system for automatically gene.rating, low-level command sequences from Ili:,h-level usergoals. 1 DCAPS uses artificial intel ligence (A1) based search techniques and an iterative repair framework in which the system selectively resolves conflicts with the resource and temporal constraints of the 1 DATA-CHASER shuttle payload activities.

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### 1. Introduction

Command sequence generation for spacecraft operations can be a laborious process requiring a great deal of specialized knowledge. Command sets can be large, with each command performing a low-level trek, There may be many interactions between the commands due to the use of resources. In addition, due to power and weight limitations, the resources available on-board spacecraft tend to be scarce. Because of this complexity, tools to assist in planning and scheduling spacecraft activities are critical to reducing the cost and effort of mission operations.

This paper describes a general system that uses Artificial Intelligence Planning and Scheduling technology to automatically generate command sequences for the DATA-CHASER shuttle payload operations. The DATA-CHASER Automated Planner/Scheduler (DCAPS) architecture presented supports (lirect, interactive commanding, rescheduling and repair, resource allocation, and constraint maintenance.

The LOCAPS search algorithm was developed based on the "iterative repair" technique used in [14]. Basically, this technique iteratively selects a schedule conflict and performs some action in an attempt to resolve the conflict. Using a repair algorithm, DCAPS is naturally well adapted for human interaction. Therefor, the scheduler can be used a s a tool to assist payload command sequencing. With the use of this tool, sequencing becomes simple enough to be accomplished by non spacecraft and sequencing experts, such as the mission scientists. This allows the scientist to become directly involved in the command sequencing process. Following any changes in spacecraft state or user-defined goals, the repair algorithm allows simple, non-disruptive reschedul ing. Finally, the highly restrictive payload resources and constraints are consistently monitored and conflicts avoided automatically.

The DCAPS system is being developed for operation of the DATA-CHASL R shuttle payload which is being managed by students and faculty of the University of Colorado at Boulder. DATA-CHASLR is a science

three solar observing instruments, Par Ultraviolet Spectrometer (PARUS), Soft X-ray/Extreme Ultraviolet Experiment (SXEE), storage, communication, and control of observation. The main activities for the payload payload, with must be sequenced while avoiding or resolving spectra. The payload resources include power, involve science instrument observations, constraints. conflicts activities can be scheduled. Payload activities orientation - which affects when certain science externally-driven states such as the shuttle instruments, and DATA-CHASER tape storage, (LASIT), that are imaging devices at various and Lyman-Alpha Solar Imaging power subsystem. Science is performed using With and the communication SFR is also constrained local memory, the ad the communication a primary focus resources and the on solar Telescope temporal three bus. data

activities. The second phase offers an interactive scheduling session. Using the repair-based scheduler, the user can work with a small set of high-level science and engineering goals, an initial schedule can be generated. The goals, which describe highconsistency with resources translated level mission objectives, are automatically modes of operation. First, by simply providing When using the DCAPS system, there are three all conflicts in the current schedule. A schedule free of conflicts, however, may not be the user can give one simple command to resolve schedule to search for a better solution. user can call on the optimizer to rate the highest quality schedule. In the final stage, the After making any change in the schedule, the low-level into a sequence The second pha activities while and constraints. of maintaining executable

The main scheduling algorithm of the planner/scheduler is the repair-based search algorithm. Using this algorithm, the scheduler first collects all of the conflicts in the current schedule and classifies them based on the resource being violated and the culprit activities associated with the conflict. After choosing a conflict to repair, the scheduler must select an action to perform in an attempt to resolve the conflict. Actions include moving, adding, and deleting activities. If the action resolves the conflict, the scheduler iterates on the resulting schedule. Otherwise, the scheduler tries a

different action for resolving the persistent conflict.

to automated command generation. Next we describe the model representation. Then we describe how DCAPS fits in to the overall command the DATA-CHASER payload. We then go into detail about the DCAPS approach objectives. Next we discuss the different ways in which the DCAPS system can be used to CHASER follows. The remainder of this paper is organized as related work and conclusions. flight and ground system architecture for the DATA-CHASER mission. Finally we discuss First shuttle we describe payload and the morsim DATA-

# 2. DATA-CHASER PAYLOAD

DATA-CHASER consists of two synergetic projects (see Fig. 1), DATA and CHASER, which will fly as a Hitchhiker payload aboard STS-85 on the International Extreme ultraviolet Hitchhiker Bridge (HEH-2) in July 1997 [3]. A technology experiment, DATA (Distribution

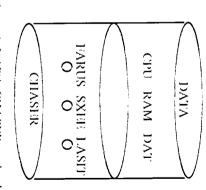


Figure 1: DATA-CHASER payload

and Automation Technology Advancement) seeks to advance semi-autonomous, supervisory operations. CHASER (Colorado Hitchhiker and Student Experiment of Solar Radiation) is a solar science experiment that serves to test DATA. The DATA technologies support cooperative operations distributed between different geographic sites as well as between humans and machines, on-board autonomy, human control, and ground automation.

CHASER is comprised of three co-aligned instruments that take data in the far and extreme

Ultra-violet wave.-lengths. The first and oldest of these instruments (17 years old) is 1 ARUS or the 1 'AR Ultraviolet Spectrometer, which takes a continuous spectrum from 115 nm to 190 nm with a resolution of .12 nm, LASIT (Lyman-Alpha Solar Imaging 'J'clc.scope.) takes images of the full solar disk of the sun in the Lyman- α wavelength (121.6 nm) with a CII) imager. The final instrument in the scientific package consists of 4 photometers, each having a different metallic coating so as to enable them to look" at different wavelengths between 1 and 40 nm. The objective of these instruments is to measure the full disk solar ultray jolet irradiance and obtain images of the sun in the Lyman-α wavelength, providing a correlation between solar activity and radiation flux as well as an association of 1 yman-α fluxes with individual active regions of the sun.

The flight segment of the 1 DATA-CHASER project consists of a canister that is equipped with a 1 litchhiker Motorized Door Assembly (HMDA) which houses the instruments and their support electronics. The second canister contains the flight computer for the payload as well as the 2 GB Digital Audio Tape (DAT) drive that is used to store all data that is collected during the mission. The payloads data is also sent to the ground system through both lowrate (available 90% of the time, @ 1200° bps) and medium rate (available when scheduled, @ 200° kbps). The payload is also capable of receiving commands sent from the ground system when uplink is available.

During the mission, the DATA-CHASER payload will be operating in four different modes. Most of the time DATA-CHASER is powered it will be in a passive mode where it is monitoring its state and notifying the ground of any changes. During the lime in the mission whenthe Orbiter is scheduled to point the bay at the sun, the DATA-CHASER payload will shift into solar active mode where each of the instruments takes data. That data is both written to the DAT drive on board and downlinked to the groundsystem for immediate data analysis. Several times during the mission, DATA-CHASER will take data while not pointing, at the sun. This data is used for testing various portions of the DATA experiment with nonsolar pointing data in addition to being used for instrument calibration.

One of the consequences of flying on the shuttle system is that shuttle resources are limited and their availability is subject to change every 12 hours. These resources include access to u plink and downlink channels, and time that your payload is allowed to operate. In addition to these resources, any given payload may also have environmental constraints as to how much contamination the payload can take. Another example is thermal constraints, such as maximum solar point time.

S'I'S- 85, the flight that D ATA-CHASER payload is scheduled to fly on, is one of the most complicated flights that the shuttle has flown to date. In addition to DATA-CHASER payload, there are 4 other payloads sharing the same H H bridge. In addition, the IEH-2 bridge there is another IIII bridge, a pallet payload, and a Spartan deployable satellite. Needless to say the shuttle pointing requirements are considerably tight. In addition to modeling what the internal constraints and resources of the payload are, DCAPS must also search the shuttle flight plan for times when we are allowed to operate, downlink our data, uplink new command set s, and when we have to protect CHASER science instruments from contamination events.

DATA-CHASER is an interesting scenario for scheduling because of the complex data and power management involved in the science gathering. An automated scheduler must find an optimal "data taking" schedule, while adhering to the resource constraints. In addition, the scientists would like to perform dynamic scheduling during the mission. As an example, the summary data may indicate the presence of a solar flare. If this occurs, scientists have different requirements and goals, such as higher priorities on certain instruments or longer integration times. These new goals may require a different schedule of activities.

### 3. USER OPERATION

The DATA-CHASER Automated Planner/Scheduler will be part of the DATA-CHASER mission operations software, It will

be a ground-based intelligent tool used for generating scheduled **commands** for uplink to the payload. The user's manual can be found at [9]. '1'here, are three phases of operating the 1 DCAPS system: a goal satisfaction phase, an interactive repair phase, and an optimization phase.

### 3.1GOAL SATISFACTION

The first phase? of scheduling in DCAPS involves generating an initial schedule from a set of high-level, user-defined goals. The scientist or engineer simply requests one or more of the goals, and the scheduler will generate the low-level activities that satisfy the goals. For example., the scientist can simply make a request for solar observations during all solar viewing periods. Given this request, the scheduler will create and position the instrument data-take activities and their supporting activities.

Goal satisfaction is a way of generating an initial schedule. Goals are parameterized, and create activities in positions relative to certain schedule events or parameters. In this way, the same goals can be requested for different initial states. This makes them more flexible than alternate ways creating an initial schedule, such as simply loading activities from a file. For example, the initial state in DATA-CHASER contains shuttle maneuver activities. These activities determine the solar viewing periods of the payload. The solar observation requests are based relative to these solar views, and therefore, are applicable to any maneuver sequence.

### 3.2 Interaction Repair

in the second phase of scheduling, the user has the opportunity for interacting with the schedule at amore detailed level. The scientist or engineer can view the activities at several levels of abstraction. The GUI can display activities from the highest level, as a single event, down to the lowest level, showing the detailed steps that make. - up the activity.

The user can also modify the schedule by moving, adding, or deleting activities, as well as changing activity parameters. For example, the scientist might want to delete a LASIT datatake and replace it with a FARUS or SXEE

(lata-take,. (h. perhaps he/she may simply want to change the tal'get of some data-take, from a solar scan to an earth scan. Although the user has the capability of making these types of adjustments, he/she does not need to worry about the various interactions, constraints, or resource usage of the activities being modified. This information is monitored by DCAPS, and changes are propagated to the dependent objects. In addition, when the user introduces scheduling conflicts, 1 DCAPS can resolve them automatically.

1 DCAPS can be called upon at any time to resolve any conflicts residing, in the schedule. Conflicts are violations of resource capacities or temporal constraints. In this way, the user does not need to be very informed, careful, or specific about his/her requests. For example, a scientist can move a (I:th-lake activity without concern for it's power usage. or, a general request for (Md-takes can be made, without specifying the exact times for the activities to occur. Although these, changes or requests may cause one or more conflicts, DCAPS can resolve these conflicts with one simple command.

### 3.3 OPTIMIZATION

1 finally, the thirdphase of DCAPS operation is schedule optimization. After resolving all conflicts, the schedule may still contain violations of user preferences. These violations can be expressed as "soft conflicts" and can be repaired in a manner similar to repairing regular conflicts. The main difference is that the modeler must explicitly represent conflicts" and general mechanisms for repairing them. As an example, considered an engineer's desire to have all dots written to permanent storage at the end of the mission, Having data in the RAM at the end of the schedule is not a violation of the resource, but could be considered a "soft conflict>".

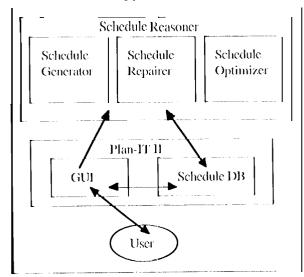
1 references can also be expressed in a schedule evaluation function. In the optimization phase, IX APS can score valid schedules based on the evaluation function developed by the modeler. This evaluation function can be used to produce more optimal valid schedules. One simple evaluation function may give higher scores to schedules with more science observations.

1 DCAPS can find more optimal schedules by running the automated scheduler many times and remembering the schedule with highest score.

### 4. Automated Planner/Scheduler

The DATA-CHASER Automated Planner/Sel reduler produces a complete, valid schedule of payload operation commands from a model, initial state., and set of high-level goals. In addition, it can input intermediate, invalid schedules (resulting from user changes) and produce a similar, but valid schedule. Finally, the scheduler can lake several valid schedules, score them, and select the most optimal schedule.

The planner/sch eduler consists of two main parts, the Plan-IT II (1'1?) sequencing tool[5] and the schedule reasoner (sec. Fig. 2). PI2 was written by William C. Eggemeyer and originally designed as an "expert resistant sequencing tool.'" 1'1? includes a Graphical User II)tel'face. (GUI) that allows for easy manipulation of the schedule. In addition, it serves as an activity/resource data-base that supplies valuable. information to the schedule reasoner. 1'12 supports complex monitoring



1 figure 2:1 )CAPS architecture

and reasoning about activities and the various constraints between them. The schedule reasoner uses Artificial Intelligence (AI) techniques to automatically generate new schedules, repair existing faulty schedules, and

optimize valid schedules. P12 provides information about resource availability and conflicts, while the scheduler must decide which activities to use to resolve the conflicts and where to place the activities temporally. The two components work together to provide easy and fast sequencing of mission activities.

### 4.1 SCHEDULE DATA-BASE

In the DCAPS system, P12 is used primarily as a "schedule (Inta-base" and resource constraint checker. It was originally (]c.vclo]k x1 as a graphical sequencing, tool. Activities and resources are displayed on a graphical output. An activity represents some mission event that occurs over a period of time and uses some of the mission resources. A resources represents some limited available material whose usage is modeled as discrete blocks over time. For each type of activity and resource, 1'12 displays a timeline, which represents the behavior of that activity/resource type over a period of time. When activities are created, they are placed at a specified time on the timeline. Resources used by that activity are updated to reflect the additional usage, i n addition to schedule visualization, Pl2 provides an easy-to-use input interface for modifying the schedule. Moving activities is as simple as a click-and-drag with a mouse.

P12 helps ease the bill'dC11 on sequencers by continually monitoring all activities in the sequence. As activities are added or moved, the change in resource usage is automatically updated, and the new resource profiles are displayed. With this in formation available, the user can immediately see the effects of a schedule change on the mission resources. For each resource, P12 also monitors any conflicts that are occurring on the resource. Conflicts are time intervals where the limitations of the resource have been exceeded. These conflict intervals are highlighted in red to flag there existence for easy identification. Finally, 1'12 monitors any dependencies that have been defined between activiti es and resources. The values of specific parameters of activities and resources may be functionally dependent on Values of other parameters. 1'12 automatically keeps these, parameter values consistent.

1'12 also helps out by serving as an activity and resource data-base, producing/accept in?, to/from sequencer. information a functional interface 101'12 has been extended to better assist an automated sequencer. A basic set of "fetch" functions have been developed to quickly retrieve information about conflicts and the resources and activities involved in the conflict. 1 or example, an interface function has been written to fetch the legal times where an activity can occur in the. schedule. 1 lere, "legal times" refers to positions where no conflicts are caused by any Of the resources used by the given activity.

In addition to fetching information droll the current stale, of the schedule, the user willneed to be able to change the current state in attempt 10 fix or optimize the schedule. Some basic printitive functions are provided by Pl2 to allow an external system to add and move activities, change there duration, etc. These primitives make up the set of actions that a scheduler can take when trying to resolve conflicts.

### 4.2 SCHEDULE REASONER

The second major component of DCAPS is the automated schedule reasoner. This is the next step in automating and simplifying the spacec raft command sequencing process. There are three parts to the schedule reasoner: schedule generator, a schedule repairer, and a schedule optimizer. First, the schedule generator will transform a set of user-(icfinc(i high-level goals into a valid sequence of lowlevel commands. Second, the schedule repairer will automatically maintain the consistency of the sequence after arbitrary user interaction by repair The rescheduling using actions. scheduler repairer iteratively attempts to resolve each conflict, which involves making choices on what to repair, and how to repair it. Finally, the schedule optimizer can optimize a valid schedule to increase science return.

### Sc hedule Generator

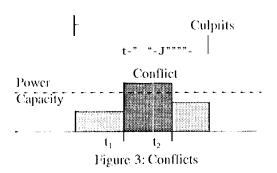
The first Step in sequencing spacecraft commands is to come up with an initial schedule of events for each phase of the mission. 'his process has been partially automated in DCAPS with the schedule generator. Expressing schedules and partial

schedules to be gen crated is done through user defined goals. There are two ways in which user goals are handled in DCAPS. First, initial science and engineering goals are addressed with parameterized scheduling functions. Each functions implements a goal. For example, there is a "Place-Power" function that schedules power switching activities in app ropriate plain based on some engineering parameters. Parameters may include s u c h things as a minimum time between switching or a power (m during a particular state of a different resource. Second, science goals can also be expressed through data-take requests, which do not have to be a part of the initial schedule generation. For example, a scientist can request ten additional scans from a particular instrument to occur any time during some phase of the mission. This type of general request does not include specific locations or necessary supporting activities. The scheduler will simply place them at random positions and allow any conflicts to be resolved by the automated repairer.

### Schedule Repairer

The generated initial schedule may still violate some of the spacecraft constraints. Also, the scientists and engineers might feel there goals were not completely satisfied, and may need to interact with and modify the generated schedule. In doing this, new conflicts may be introduced. Therefor, we need some way O f automatically resolving any existing conflicts in the schedule, while (disrupting the currentstate of the schedule as little as possible. Having the p rocess automated allows the user to be less careful, and therefor spend less time. on the details of sequencing the activities. When general requests 01" changes have been made, all conflicts can be resolve. (i by executing one simple command to invoke the schedule repairer.

Before describing the schedule repairer, we must present a few definitions. A "hard conflict", or just "conflict", is a violation of one of the resource constraints. A conflict occurs over a certain time period and is caused by activities called "culprits". For example, if the power capacity is exceeded from time  $t_1$  to time  $t_2$ , then there is a conflict from time  $t_1$  to time  $t_2$  and the culprits are any activities that use power



(luring this time (see Fig. 3). A "soft conflict" is aviolation of one of the user's high lewd goals. Hard conflicts are violations" of legal constraints, while. "soft conflicts" are violations of user preferences. "Choice points" are places in the scheduling algorithm when a decision must be made. For example, when there are many conflicts to resolve, the scheduler must decide which conflict to resolve first. A "hard choice", or just "choice", is a decision made solely on the basis of possible hard conflicts. It may be decided, for example, not to place an activity at a certain time bc.cause. new conflicts will be added as a result of that placement. A "soft choice" is a decision made on the basis of user preferences or a heuristics with the hopes of generating a more optimal schedule. An example of a user preference is a priority scheme 01) certain activities. One heuristic may be to move lowest priority culprits 10 the nearestlegal position.

There are four possible actions to take in attempt to resolve. a conflict: move, add, or delete an activity. The "move." action involves moving (me. of the culprits of the conflict to a positions that will either resolve the conflict, or at least insure that the moved activity is no longer a culprit. Some conflicts can be resolved by adding a new activity. These activities usually provide some resource that was previously not available. Finally, a conflict can also be resolved by simply deleting the culprits. This is obviously not a preferred method and is Only used as alastresort.

The resolution of a conflict greatly depends on the type of resource that is in violation. There are five different types of resources, and therefor, five different types of conflicts to resolve. A conflict on a depletable resource means that the activities of the schedule have used too much of the resource. In this type of conflict, the culprit is the activity that caused the resource to overflow" during the time that it first overflows. Some depletable resources have "resetter" activities and this sort of conflict can be resolved by adding an activity that "resets" the available resource. For example, a downlink activity will free up spare in the downlink buffer. A conflict on a non-depletable resource is when activities o\'c.ruse a resource during a particular time interval. The culprits in this type of conflict are all of the activities that use the resource during the conflict interval. This sort of conflict can be resolved by moving or deleting culprits. There are no activities in the DATA-CHASER model that can add to a non-depletable resource. A conflict on a state resource is when an activity requires the resource to be in a state which it is not. The culprits in this type of conflict are all of the activities that require. the incorrect Wile, and the activity that changed the eresource to the incorrect state. Several possibilities for resolving a state conflict include moving the culprits to another interval where the required state is present, or adding an activity that will change the state of the resource to the required state. A conflict on a concurrency resource is when an activity requires the presence of the resource, which is usually provided by the another activity. The culprits in this type of conflict are all of the activities that require the presence of the resource. To resolve concurrency conflict, the scheduler can move the culprits to an interval where the resource is present, or add an activity that provides the presence of the resource. Anti-Concurrency conflicts are essentially the same except they require the absence of the resource rather than the presence of it.

For any type, of initial schedule, the schedule repairer must find the correct activities to move, add, or delete and position" them temporally in such a way that no conflicts remain. The scheduler is based on a random scheduler with several heuristics used at the various choice points. The scheduler relies (m some interface functions to 1'12 that describe the conflicts in the current schedule, describe the activities that could resolve a conflict, and manipulate the schedule. We first describe the random scheduler, followed by the. heuristic enhancements that facilitate scheduling within the DATA-CHASER domain. The ultimate task

The following is the algorithm for the schedule repairer written in a C-like psedo-code."

```
Resolve-Conflicts ()
   iterations = 1
   conflicts : GetConflicts()
   hoop while (length(conflicts) > 0 && iterations <= max-iterations) {</pre>
      conflict = ChooseConflict(conflicts)
      method = ChooseMethod(conflict)
      case (method) {
      'move'
         culprit : ChooseCulpritToMove(conflict)
         duration = ChooseDuration(conflict, culprit)
         start-time = ChooseStartTime(conflict,culprit,duration)
         success = MoveCulprit(conflict,culprit,start-time)
      'add'
         activity = ChooseActivityToAdd(conflict)
         duration = ChooseDuration(conflict, activity)
         start-time = ChooseStartTime(conflict,activity,duration)
         success : AddActivity(conflict,activity,start-time)
      'delete'
         culprit = ChooseCulpritToDelete(conflict)
         success : DeleteCulprit(conflict,culprit)
      progress : GetProgress()
      if not (success && progress) then UndobastAction()
      conflicts : GetConflicts()
      iterations : iterations + 1
```

of the system is to find the best place to schedule the activities so as to maximize the utility of the schedule. In the basic random scheduler, all choices are made randomly from the list of **options** unless otherwise specified. The algorithm is a simple iterative. loop over the conflicts in the schedule. First, a conflict is selected from the list of current conflicts. An attempt is made to resolve the chose.n conflict. Next, a method for resolving the conflict is chosen. The repair action will depend (m which method has been selected. If "move" is chosen, then a culprit must be picked from the list of culprits in the conflict. A duration and start time are chosen for the. culprit, and the culprit is moved to the new lo cation. If "add" is the chosen method, then the repairer must decide which activity type to instantiate. Again, a duration and start time must be chosen for the new activity, and the activity is inserted at the chosen time. If the repairer chooses to "delete"

an activity, then it simply must choose an activity to delete, and delete it. After the chosen action is performed, the schedule repairer clinks to see if progress was made. If the action did not succeed in resolving, the conflict, or progress was not made, then the action is "undone". Otherwise, the new set of conflicts are found, and the loop counter is incremented. This process continues until all conflicts are resolved, 01 the loop counter exceeds a user defined maximum bound. 1 for every choice point in the algorithm, where a selection must be made from a list of possibilities, the schedule repairer is all lewd to "backtrack" to that point. What Ibis means is, that if a particular choice fails, it may choose another from the list before giving up. If all choices fail, then a previous decision must have been incorrect, and the repairer can "back track" to the preceding choice point. All choice points, including the decision on whether or not to backtrack, are heuristic decisions, and may customized to a particular domain.

### Schedule Optimizer

There are three ways to optimize a schedule: using preference heuristics at scare.h choice points, specifying a set of "soft conflicts", or by simply scoring results from multiple. runs of the scheduler. A preference heuristic, 01" "soft choice", can be made at any decision in the repair search. 1 for example, when decid i ng where to move. a conflict causing activity, the user might prefer to move that activity to a position close.st to it's current position. This will help the scheduler avoid unnecessary disruption to the existing schedule. The existing schedule, after all, may have been produced by the user in an attempt to optimize the schedule.

1 Preferences can be expressed using what we referred to as "soft conflicts". A soft conflict is a way of specifying a preferred value for a particular resource, possibly at a particular time. I 'or example, having any scanned data that has not been stored on the tape at the end of the mission, is considered a soft conflict, The scientist would prefer that all of the data be written to the tape at the missions end, rather than leaving, it in the on-board memory. After the schedule repairer handles all of the "hard conflicts", it continues by iteratively addressing all of the "soft conflicts".

The third approach to optimization involves scoring several resulting schedules and choosing the one with the highest score. The evaluation function is domain dependent, and would have to be written separately for each application. 1 lowever, some basic scoring will be similar across applications. 1 for example, most science spacecraft are mainly concerned with collect the laurgest number of images as possible. A simple evaluation would give a higher score to schedules with greater amounts of collected dat a, Once we have the evaluation function, we need to be able to produce several (Iirfc.rent schedules from the same goals and initial stale. This can be (1 on e by either changing t h e heuristics, or running the sc] icduler with a different random seed. Some heuristics may work better than others, and it is often difficult to tell which is the best for a

particular application. '1 'herefore, it may be necessary to resort to empirical tests. After running the scheduler on different heuristics. we can simply choose the set of heuristics which generates the schedule with the highest Score. After choosing the heuristics, the scheduler can be run many times with different random seeds. At choice points where there, is no heuristic for choosing from the list of possibilities, the scheduler makes a random decision. With different random seeds, these decisions will be different, and the resulting schedule will be different. Using the evaluation function, we can assign a score to each, and choose the schedule with the highest score. This procedure will not necessarily uncover the optimum schedule, but it will help find a more optimal schedule.

### Heuristics

'I he general search and decision making described above, would be futile without expert support and guidance. I leuristics have been developed and i neorporated into 1 DCAPS to help guide the, search to a valid and more optimal schedule. This guidance knowledge comes from both domain experts and scheduling experts, '1 here are three basic classes of heuristics used in DCAPS: selection, pruning, and backtracking heuristics.

Selection heuristics involve deterministical I y sorting or selecting from a list of possibilities at a choice point in the search, The selection is usually based on some property Of the objects being considered. For example, when choosing a culprit to move in order to resolve a power conflict, one heuristic might choose the culprit that uses the most amount of power. Using this heuristic might resolve the conflict faster. Another successful heuristic used in DCAPS was one that sorted the possible locations for act ivity placement by the. number of conflicts the activity would cause when placed in that location, 'his basic approach has been referred to as the "min-conflicts" heuristic [7]. The minconflicts algorithm we use is interesting and worth going into detail.

I 'or each resource used by an activity, we query the (i at a-base for the legal times where the activity can be placed without violating the resource constraint. Then, each legal interval is

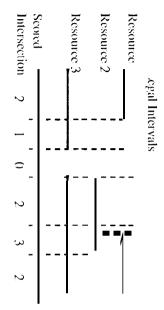


Figure 4. Min-conflicts with scored interval intersection

result of this intersection is then intersected no interval where neither A nor B exist. The two of the resources, using a special "scored" with the fewest conflicts. the highest score, in other words, the position the activity is placed in that position. Using these intervals we can choose a position with number of resources that will not be violated if intervals, been intersected. The result is a set of scored until each set of intervals for each resource has with the third set of intervals. This continues score of B where the two intervals intersect, or does not, an interval with a score of A plus the interval with a score of B where B exists and A possibilities: an interval with a score of A for intersection of intervals A and B results in four interval intersection (see Fig. intersect two sets of intervals that resulted from assigned an initial score of one. positions where A exists and B does not, where the score represents 4). The scored Next, WC Ê

be to prune only those intervals that would schedule. A more conservative approach might have included positions that, if the activity was possibilities that could be useful. In the above futile and giving up to try something different. few positions before realizing this attempt is scheduling because the scheduler will only try a number of conflicts. try positions with the highest score, or least possible positions. One possibility is to only for an activity, we may not want to try all For example, after finding the scored intervals attempt to make the choice easier and faster. some of the possibilities for a given selection in the pruning heuristics. These heuristics remove Another class of heuristics used in DCAPS are example, some of the pruned intervals may However, too much pruning may remove there, would have This may speed up improved

cause more conflicts than are currently in the schedule. These intervals cannot possibly be positions that could improve the schedule.

may fix the chosen conflict, but cause several other conflicts. Therefore, success can be requires an approximate definition. Success is the action, or try a different one. The second Checking the progress of an action can be used as a heuristic for determining whether to accept and the change in the duration of the conflicts. conflicts, the change in the number of culprits, much? And what defines an improvement? Our thought of as improving the schedule. But how notion of an "action failure" is not clear and backtracking: deciding on "action failures" and choice point. Heuristics can be used to help the next possibility, or move on to a different same problem and when to move on to a determine when to continue working on the previous choice was a failure. Heuristics can point, at some point we must decide that the and failing on, a list of possibilities for a choice opportunity for heuristics comes when deciding if there is a "selection failure". While trying, obscrving current When, resolving a conflict, and action attempt deciding on possibility, and it fails, we can continue and try have a list of different problem. At each choice point, we Finally, backtracking heuristics are used to help help with this decision also. simply resolving OW1 definition definition of progress inclusion the change in the number "selection failures". First, types of possibilities. If we try the decisions chosen conflict. includes

# 5. Modial Representation

In order to use either Plan-IT II stand-alone or the full DCAPS system, the user must write a software model of the mission activities and spacecraft resources. This involves defining a set of objects and how they interact. These definitions are then used by the scheduler to create instances of the objects.

### 5 MODELOBJECTS

The basic objects in the P12 sequencing tool are activities, resources, slots, and dependencies.

Activities

Activities are used to model the events that happen that affect the DATA-CI IASI R payload, and the actions that the DATA-CHASI R payload can take. All activities have some basic components: a duration, a list of slots, and a list of slot-value assignments. In addition, certaintypes (described below) have a list of sub-activities. For the sc activities, the user can also define a set of temporal constraints between the sub-activities. Next, we describe in more detail the four basic types of activities: events, steps, step-activities, and activities.

Events are used to model activities that do not occur in a fixed relation to other activities (like TDRSS contacts) and are not part of an activity hierarchy.

Steps are the "hi" nodes in the activity hierarchy tree. In other words, they do not contain any sub-activities. Steps cannot be instantiated without their parents and arcused to mode! the activities at the lowest level of detail. For instance, we model an activity called CHASER-heating, which consists of two steps, CHASER-heater-on and CHASER-heater-off.

Step-activities are used to model activities at a middle level of abstraction. They can contain steps, but must also have parent activities. In DCAPS, we model an activity SXEE-I DataTake, which models the SXEE instrument opening it's aperture and taking a scan. In this case, there is a step-activity called SXEE-ScanStep, which has sensor read steps and can not be instantiated by itself.

Activities are used to model activities at the highest level of abstraction. They are the "root" nodes in the hierarchy tree, containing subactivities, but no parent activity. The activity and event objects are what the scheduler can instantiate, and Plan-It II provides methods to access the varying levels of abstraction.

### Resources

R esources define the various physical resources and the constraints they impose. Resources come in essentially four varieties: state, concurrency, depletable, non-depletable, and simple..

Stateresources are used to 111(XIC1 the systems in the DATA-CHASER payload which have states associated with them. 1 for each state resource, the modeler must specify the possible values that the state can be. Most of the systems have at least one state variable, which is whether or not they are activated. The orientation of the payload is also modeled with a state variable which is discretized into four Slates (solar, lunar, Cal'[ii, deep space.).

Concurrency resource const raints are used to model rules that stipulate that an activity either must occur with another activity or can not occur with another activity. One relationship that is modeled with a concurrency resource is the requirement that a downlink or uplink can only occur during contact with a TI DRSS satellite. This is modeled as a resource that is present when there is TDRSS contact activity, and required when there is a downlink or uplink activity.

1 Depletable resources are used to model resources with a fixed quantity, such a fuel or RAM, Activities can List some. finite amount of a depletable resource, which may or may not be restorable. The amount used by the activity is persistent to the end of the schedule. In addition, the modeler must specify a maximum capacity for each depletable resource. In IDC APS, RAM is modeled as a depletable resource. Science observations produce data and use some amount. Of the depletable resource. Other activities, such as a transfer to permanent storage, may restore this resource.

Non-depletable resources are used to model resources which have a limit to the usage at any one time, but are reset at the end of the activity which **consumes the** resource. Similar to depletable resources, non-depletables are assigned a maximum capacity. Resources like power are modeled with non-depletable resources.

Simple resources are used to model devices which can only be used by one activity at a time. For instance, each of the instruments on board DATA-CHASER, FARUS, SXEE, and LASIT, are capable of taking only one image at a time, and are modeled with simple resources.

Simple resources are essentially non-depletable resources with an capacity of (me.

### Slots

Slots are parameters of activities that allow them to affect resources. They are defined separately, but referenced inside a ctivity definitions along with a value assignment for each slot. In the slot definition, the modeler must specify which resource it affects. The main types of slots are: info slots, simple slots, availability slots, choice slots, amount slots, and slate slots.

Info slots are for embedding text information in activities. They are merely placeholders and (10 not have any effect on scheduling.

Simple slots are included in activity type definitions in order to model usage of a simple resource. For instance, there is a slot called FARUS which is included in activity definitions of activities which use the FARUS instrument. This is how usage, of the FARUS instrument is modeled.

Availability \$10 ts are the stots that allow activities to provide or require the presence of a concurrency resource. There is a slot in 1 DCAPS called TDRSS-coverage and a slot called '1'1 DRSS-coverage-needed. Both affect the '1'DRSS-coverage resource. '1'1)1<ss activities have the TDRSS-coverage slot, and downlink activities have the. TDRSS-coveragenceded slot. TDRSS activities can be placed anywhere, and provide the presence of the resource. Downlinks can only be placed in intervals where TDRSS activities are. present, since this activity possesses the slot which requires the '1'1 DRSS resource to be present.

Amount slots come in two varieties: amount and reset-amount. Amount slots reduce a depletable of 1 non-depletable resource, and reset-amount slots increase a depletable or non-depletable resource. Amount slots do not have to be associated with a resource, however. In DC APS, w.c. have an amount slot called Rate, which is how we model the different bit transfer rates inactivities that move data, such as a down inks or DAT reads and writes. To find the amount of data an activity transfers, we multiply the rate by the duration of the activity.

There are also two types of state slots: state-users and state-changers. State-usel" slots require the presence of a certain state in a state resource, and state-changer slots change the state of a state resource. The modeler must define the set of possible states. In DCAPS, there is a state resource that models the shuttle orientation, which can be solar, earth, lunar, or deep-space. Solar science activities require the shuttle orientation State to be solar, while shuttle maneuver activities change the orientation state.

### **Dependencies**

Plan-Itll provides the ability to set uplinks that all low one object 10 affect another object. These links are called dependencies. There are several 1 ylpes of dependencies based on the types of objects it relates: slot to resource, slot to slot, slot to activity startor duration, activity startor duration to slot, and resource to resource dependencies.

Slot to resource dependencies are the default dependencies in the Plan-It 1 I system. They allow a slot to affect a resource, and are created automatically when a slot is defined with the same name as a resource.

Slotto slot dependencies allow the value of one slot to affect the value of another slot, 1 for instance, in the 1 DAT-transfer activity, there are two slots, one that models the removal of data from the RAM, find one that models the addition of data to the DAT (digital tape). In DCAPS, a dependency has been defined that sets the value of one of the slots equal to the value of the 0111(V" slot (so that the amount subtracted from RAM is never different than the data added to the DAT).

Activity start time or duration to slot dependencies and slot to activity start time or duration dependencies facilitate the modeling of convenient relationships among Plan-It II objects. For instance, the DAT-transfer has a slot called Rate, which is the rate at which data can be moved from the RAM to the DAT. We have a dependency which sets the amount of data that is removed from the RAM equal to the rate multiplied by the duration. An example of a dependency which goes from slot to duration is

a dependency which links the selected target for a science image to the length of time it takes for the instrument to scan. The duration of a 1 ARUS scan varies depending on it's use of the shuttle orientation state (solar, earth, or lunar).

Resource to resource dependencies allow one resource to affect another resource (1 i rect ly. This is very convenient for modeling power usage, since power consumption can be tied to a ctivities or stale.s. For instance, power consumption by the heater can be linked to an activity (e.g. the activation of the heater), or to a state of the heater (e.g. When the Stale, of the heater is "m", more power is use(1).

### 5.2. HIERARCHY

The modeler can create an activity hierarchy when defining the activities. All this means is that activities can have sub-activities which can also have sub-activities, and so on. Only the activity at the. top of the hierarchy can be instantiated in the schedule. When an activity is created, all of it's sub-activities are created automatically. Therefore, the scheduler must schedule the entire hierarchy if it wants to schedule one of the components.

In modeling the. DATA-CHASER shuttle payload, decisions had to be made about where to put activities in the activity hierarchy. We decided to model those activities which could be scheduled arbitrarily (and had no subactivities) as events not in a hierarchy. Some activities that were modeled as events were TDRSS contacts, shuttle venting, and very simple activities which (X)111(1 occur independently, like relay activations and 11 MDA operations (opening and closing).

If one event always occurred in some fixed ten iporal relationship to another, then we modeled it as an activity in a hierarchy. For instance, a SX EE data take consists of a number of calibrations, a door opening activity, several scans, a door closing activity, then a data transfer to the RAM buffer. We modeled all of these activities as steps in an activity called SX EE-Data-Take.

### 5.3 COMMON STRATEGIES

There were a number **of** strategies that we employed in the modeling process that m a de modeling the DATA-CHASER payload simpler.

One strategy that we employed was to reduce the number of states that state variable.s could have through discretization. For instance, spacecraft orient alion is best modeled with a real valued 3 dimensional vector. But for modeling purposes, we reduced the number of possible orientations to a discrete set of four: solar, lunar, earth, and deep space.

Another strategy that we employed in modeling D ATA-CHASI R was to separate one. component into several. I for instance, there was really Only one memory buffer that was u sc(1 for storing several types of data, but we modeled it as though it were three buffers: one for science data, one for engineering data, and one for storing data to be downlinked. We also did this with power. There are really only two power sources, IDATA power and CHASI R power, but we modeled them as though there were d i fferent power resources for each of the science instruments and several of the. other subsystems. This allowed us to track power usage more conveniently.

### 6. System Integration

1) CAPS will be integrated into the 1 ind-to 1 ind Mission Operations System (EEMOS) that is current by being developed for the DATA-CLI ASI is project as a prototype for the 1 luto Express EEMOS [1] O]. Currently the 1 DATA-CHASER EEMOS consists of 7 palls: Command & Control, 1 ault/1 ivent Detection InteractionReaction (1 / EDIR), 1 DATA/IO (1 Data handling), the Ground Database, the Graphical User Interface, the software testbed, and finally the planning and scheduling system (DCAPS).

The command and control system that we are using System Command 1 anguage (SCL, also known as Spacecraft Command Language) which integrates procedural programming with a wal-time forward chaining rule based system. 1 DCAPS interfaces with SO, through DATA/IO by sending script scheduling commands that are to be scheduled either on the flight or ground system. This is (lone by mapping PI2 activities to SCL scripts that were written prior

to flight. These commands are the.n Sent to 1DATA/IO where it forwards that list to the SCL. Compiler. once compiled they are sent to the payload through the uplink.

1 DCAPS is also interfaced with the ground EEMOS database, O2. O2 is an object oriented data base that will be used to store all mission data and telemetry that is downlinked by the payload. It will also store a command history. Through DATA/IO, DCAPS will request current payload status data in the form of sensor values in the telemetry history. It will also request lists of all commands uplinked during a given time interval. These are used by DCAPS to infer command completion status as well as to get the current state of the payload so that a new schedule can be created.

During mission operations, approximately every 4 hours or so, DCAPS will be asked by an operator 10 generate script scheduling commands and rule activations for the. next 6 hours according to its schedule. Once this list is finished, it is reviewed by the Mission Operations staff on duty. If judged to be cwrc.et, scheduled scripts commands will be sent to DATA/IO during the next available uplink window.

If during that 6 hour period there is a major change in the NASA activities, DCAPS will ask if the users want to update the scheduled scripts on-board. Due to the fact that SCL currently has no scheduled script instance identification, this will involve (Ic. scheduling ALL remaining scripts in the queue and then rescheduling them. This is fine if the user did not schedule any scripts independently of DCAPS. If he did, and DCAPS reschedules its list, the user's scheduled commands will be 10SI. If the user accepts it, DCAPS will generate a updated list, ask the users to verify it, and then deschedule rest of the 01(1 list and schedule the new list.

### 7. SUMMARY AND RELATED WORK

Iterative algorithms have been applied to a wide range of computer science problems such as traveling salesman [8] as well as Artificial Intelligence Planning [2,6,1 1,1 3]. Iterative repair algorithms have also been used for a number of scheduling Systems. The

GERRY/G PSS system [14] uses iterative repair with a global evaluation function and simulated annealing to schedule spat.c shuttle ground processing activities. The Operations Mission Planner (OMP) [1] system used iterative repair in combination with a historical model of the scheduler actions (called chronologies) to avoid cycling and getting caughtin local minima. Work by Johnston and Minton [7] shows I10W the min-conflicts heuristic can be used not only for scheduling but for a wide range of constraint satisfaction problems. The 01'1s system [12] can also be performing iterative repair. However, OPIS is more informed in the application of its repair methods in that it applies a set of analysis measures to classify the bottleneck before selecting a repair method.

In summary, DCAPS represents a significant advance from several perspectives. First, from a mission operations perspective, DCAPS is important in that it significantly reduces the an jount of effort and knowledge required to generate command sequences to achieve mission operations goals. Second, from the standpoint of Artificial 1 ntelligence applications, DCAPS represents a significant application of planning and scheduling technology to the complex, real-world problem of spacecraft commanding. Third, from the standpoint of Artificial Intelligence Research, 1 DCAPS mixed initiative approach to initial schedule generation, iterative repair, and schedule optimization represents a novel approach to solving complex planning and scheduling problems.

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